

DS-210: PROGRAMMING FOR DATA SCIENCE

LECTURE 39

- 1. COMPILING RUST TO WEBASSEMBLY
- 2. BONUS: FAST ALGORITHM FOR COMPUTING FIBONACCI NUMBERS



OFFICE HOURS

• Today: 4:30-6 (will finish slightly early)



OFFICE HOURS

- Today: 4:30-6 (will finish slightly early)
- Proposed final project coding jams:
 - Today: 7-???
 - Thursday: 8-???
 - (and another one on Monday?)





[REVIEW OF SOLUTIONS TO HOMEWORKS 9 & 10]



1. COMPILING RUST TO WEBASSEMBLY

2. BONUS: FAST ALGORITHM FOR COMPUTING FIBONACCI NUMBERS



RUNNING ARBITRARY CODE IN A BROWSER

Traditional way:

- JavaScript
- Java
- Adobe Flash

Shortcomings (various degrees):

- far from native speed
- portability
- many other



WEBASSEMBLY

- WebAsembly

 = "assembler" for the web
- Near-native speed
- Clear definition
- Every major web browser can run it
- Can be compiled to from many languages, including Rust



WEBASSEMBLY

- WebAsembly

 = "assembler" for the web
- Near-native speed
- Clear definition
- Every major web browser can run it
- Can be compiled to from many languages, including Rust

Additional features:

- memory directly accessible to JavaScript (to avoid foreign function interface translation)
- no garbage collection
 - adding it is being considered
 - cannot be compiled directly into from languages that use garbage collection





HOW DOES ONE USE WEBASSEMBLY IN PRACTICE?

- It does not replace JavaScript completely
- Find out which parts of your code are slow
- Rewrite them in Rust and compile to WebAssembly
- Use JavaScript on your webpage to interact with WebAssembly binaries





HOW TO DEPLOY IT

Some basics:

- Write a library in Rust
- Compile as a dynamic library
- Set the compilation target ("architecture") to wasm32 unknown unknown
- Library binaries: .wasm
- See the tutorial how to make Rust code and JavaScript talk to each other

Tutorial: https://rustwasm.github.io/docs/book/





- 1. COMPILING RUST TO WEBASSEMBLY
- 2. BONUS: FAST ALGORITHM FOR COMPUTING FIBONACCI NUMBERS



FIBONACCI NUMBERS: ALGORITHMS WE SAW SO FAR

$$F_k = \begin{cases} 0 & \text{if } k = 0 \\ 1 & \text{if } k = 1 \\ F_{k-2} + F_{k-1} & \text{if } k > 1 \end{cases}$$

Assuming O(1) time arithmetic operations:

- $O(F_k)$ time by directly recursively following the definition
- ullet O(k) time by storing values and computing F_k from F_{k-2} and F_{k-1}



Useful matrix

$$A = \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix}$$



Useful matrix

$$A = \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix}$$

Observation:

$$A \begin{bmatrix} F_k \\ F_{k-1} \end{bmatrix} = \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix} \begin{bmatrix} F_k \\ F_{k-1} \end{bmatrix} = \begin{bmatrix} F_{k+1} \\ F_k \end{bmatrix}$$



Useful matrix

$$A = \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix}$$

Observation:

$$A \begin{bmatrix} F_k \\ F_{k-1} \end{bmatrix} = \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix} \begin{bmatrix} F_k \\ F_{k-1} \end{bmatrix} = \begin{bmatrix} F_{k+1} \\ F_k \end{bmatrix}$$

Hence, by induction:

$$A^{k} \begin{bmatrix} 1 \\ 0 \end{bmatrix} = A^{k} \begin{bmatrix} F_{1} \\ F_{0} \end{bmatrix} = \begin{bmatrix} F_{k+1} \\ F_{k} \end{bmatrix}$$



Useful matrix

$$A = \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix}$$

Observation:

$$A\begin{bmatrix} F_k \\ F_{k-1} \end{bmatrix} = \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix} \begin{bmatrix} F_k \\ F_{k-1} \end{bmatrix} = \begin{bmatrix} F_{k+1} \\ F_k \end{bmatrix}$$

$$A^k \begin{bmatrix} 1 \\ 0 \end{bmatrix} = A^k \begin{bmatrix} F_1 \\ F_0 \end{bmatrix} = \begin{bmatrix} F_{k+1} \\ F_k \end{bmatrix}$$

Hence, by induction:

$$A^{k} \begin{bmatrix} 1 \\ 0 \end{bmatrix} = A^{k} \begin{bmatrix} F_{1} \\ F_{0} \end{bmatrix} = \begin{bmatrix} F_{k+1} \\ F_{k} \end{bmatrix}$$

Can we compute A^k efficiently?



EXPONENTIATION BY SQUARING

Halving the exponent:

Even *k*:

- Compute recursively $A_{\star} = A^{k/2}$
- Return A^2_{\star}

Odd k:

- Compute recursively $A_{\star} = A^{(k-1)/2}$
- Return $A^2_{\star} \times A$



EXPONENTIATION BY SQUARING

Halving the exponent:

Even *k*:

- Compute recursively $A_{\star} = A^{k/2}$
- Return A_{\star}^2

Total: $O(\log k)$ arithmetic operations

Odd k:

- Compute recursively $A_{\star} = A^{(k-1)/2}$
- Return $A^2_{\star} \times A$



EXPONENTIATION BY SQUARING

Halving the exponent:

Even *k*:

- Compute recursively $A_{\star} = A^{k/2}$
- Return A_{\star}^2

Total: $O(\log k)$ arithmetic operations

Odd k:

- Compute recursively $A_{\star} = A^{(k-1)/2}$
- Return $A^2_{\star} \times A$

Let's implement it!



AVOIDING BIG NUMBERS

- To avoid dealing with big numbers, let's just compute it modulo a large prime
- First, we have to find a big prime





AVOIDING BIG NUMBERS

- To avoid dealing with big numbers, let's just compute it modulo a large prime
- First, we have to find a big prime



AVOIDING BIG NUMBERS

- To avoid dealing with big numbers, let's just compute it modulo a large prime
- First, we have to find a big prime





"SLOW" O(k)-TIME IMPLEMENTATION

```
In [4]: fn fib_mod_linear(x: u128) -> u128 {
    if x == 0 {
        return 0;
    }
    let mut y = 1;
    let mut fib_prev = 0;
    let mut fib = 1;
    // invariant:
    // * fib_prev == F(y - 1) mod BIG_PRIME
    // * fib == F(y) mod BIG_Prime
    while y < x {
        y += 1;
        (fib_prev,fib) = (fib,(fib_prev+fib) % BIG_PRIME)
    }
    return fib
}</pre>
```



MATRIX OPERATIONS

```
In [5]: // Matrix shape
// 0 1
// 2 3

type MyMatrix = [u128;4];

const A: MyMatrix = [1,1,1,0];

fn multiply(x: MyMatrix, y: MyMatrix) -> MyMatrix {
    let mut solution = [0;4];
    solution[0] = x[0] * y[0] + x[1] * y[2];
    solution[1] = x[0] * y[1] + x[1] * y[3];
    solution[2] = x[2] * y[0] + x[3] * y[2];
    solution[3] = x[2] * y[1] + x[3] * y[3];
    solution.iter_mut().for_each(|x| *x = *x % BIG_PRIME);
    solution
```



IMPLEMENTATION OF THE FAST ALGORITHM

```
In [6]:

// exponentiation of A by squaring (module BIG_PRIME)
fn exponentiate_fib_matrix(exponent: u128) -> MyMatrix {
    if exponent == 0 {
        return [1,0,0,1];
    }
    let tmp = exponentiate_fib_matrix(exponent / 2);
    if exponent % 2 == 0 {
        multiply(tmp, tmp)
    } else {
        multiply(multiply(tmp, tmp), A)
    }
}
```



IMPLEMENTATION OF THE FAST ALGORITHM

```
In [6]: // exponentiation of A by squaring (module BIG_PRIME)
        fn exponentiate_fib_matrix(exponent: u128) -> MyMatrix {
            if exponent == 0 {
                return [1,0,0,1];
            let tmp = exponentiate_fib_matrix(exponent / 2);
            if exponent % 2 == 0 {
                multiply(tmp, tmp)
           } else {
                multiply(multiply(tmp, tmp), A)
In [7]: // Fibonacci computation
        fn fib mod logarithmic(x: u128) -> u128 {
            if x == 0 {
                0
            } else {
               let matrix = exponentiate_fib_matrix(x - 1);
                matrix[0]
```



```
In [8]:
use std::time::SystemTime;
// see how long something is executing
fn time_it(f: impl FnOnce() -> u128) {
    let before = SystemTime::now();
    let result = f();
    let after = SystemTime::now();
    println!("Time: {:.3?}", after.duration_since(before).unwrap());
    println!("Computed number: {}\n", result);
}
```



```
In [8]: use std::time::SystemTime;
         // see how long something is executing
         fn time_it(f: impl FnOnce() -> u128) {
             let before = SystemTime::now();
             let result = f();
             let after = SystemTime::now();
             println!("Time: {:.3?}", after.duration_since(before).unwrap());
             println!("Computed number: {}\n", result);
In [10]: let k: u128 = 10;
         time_it(|| fib_mod_linear(k));
         time_it(|| fib_mod_logarithmic(k));
         Time: 970.000ns
         Computed number: 55
         Time: 1.651µs
         Computed number: 55
```



```
In [8]: use std::time::SystemTime;
         // see how long something is executing
         fn time_it(f: impl FnOnce() -> u128) {
             let before = SystemTime::now();
            let result = f();
             let after = SystemTime::now();
             println!("Time: {:.3?}", after.duration_since(before).unwrap());
             println!("Computed number: {}\n", result);
In [11]: let k: u128 = 1000;
         time_it(|| fib_mod_linear(k));
         time_it(|| fib_mod_logarithmic(k));
         Time: 13.500µs
         Computed number: 202736284353
         Time: 3.105µs
         Computed number: 202736284353
```



```
In [8]: use std::time::SystemTime;
         // see how long something is executing
         fn time_it(f: impl FnOnce() -> u128) {
             let before = SystemTime::now();
            let result = f();
             let after = SystemTime::now();
             println!("Time: {:.3?}", after.duration_since(before).unwrap());
             println!("Computed number: {}\n", result);
In [14]: let k: u128 = 1 000 000;
         time_it(|| fib_mod_linear(k));
         time_it(|| fib_mod_logarithmic(k));
         Time: 12.549ms
         Computed number: 863350906745
         Time: 5.494µs
         Computed number: 863350906745
```



```
In [8]: use std::time::SystemTime;
         // see how long something is executing
         fn time_it(f: impl FnOnce() -> u128) {
             let before = SystemTime::now();
             let result = f();
             let after = SystemTime::now();
             println!("Time: {:.3?}", after.duration_since(before).unwrap());
             println!("Computed number: {}\n", result);
In [16]: let k: u128 = 1 000 000 000;
         time_it(|| fib_mod_linear(k));
         time_it(|| fib_mod_logarithmic(k));
         Time: 12.108s
         Computed number: 129171585224
         Time: 8.383µs
         Computed number: 129171585224
```